

(19)

Europäisches Patentamt

European Patent Office

Office européen des brevets



(11)

EP 1 339 198 A1

(12)

EUROPEAN PATENT APPLICATION

(43) Date of publication:
27.08.2003 Bulletin 2003/35(51) Int Cl.7: H04L 29/06, H04L 12/46,
H04J 3/16

(21) Application number: 02290445.2

(22) Date of filing: 22.02.2002

(84) Designated Contracting States:
AT BE CH CY DE DK ES FI FR GB GR IE IT LI LU
MC NL PT SE TRDesignated Extension States:
AL LT LV MK RO SI(71) Applicant: ALCATEL
75008 Paris (FR)

- Pessina, Alberto
20052 Monza (Milano) (IT)
- Pozzoli, Elena
20034 Giussano (Milano) (IT)
- Sorbara, Giuseppe
20031 Cesano Maderno (Milano) (IT)
- Ghisio, Marco
20010 Casorezzo (Milano) (IT)

(72) Inventors:

- Maggio, Santo
20064 Gorgonzola (Milano) (IT)
- Panonzini, Massimo
20059 Vimercate (Milano) (IT)

(74) Representative: Lamoureux, Bernard et al
Alcatel
Intellectual Property Department,
Vimercate
Via Trento, 30
20059 Vimercate (IT)

(54) Enhanced transport of ethernet traffic over a transport SDH/SONET network

(57) Disclosed is a method and device for handling Ethernet frame signals in a SDH/SONET network, the SDH/SONET network comprising network elements or nodes and fiber connections connecting the network elements, the method being characterized by the step of defining a new layer/network over the SDH/SONET network in order to manage the Ethernet signals over the SDH/SONET network, the new layer/network using the resources of SDH/SONET network in such a way as to optimize the provided services and the performances with reference to this specific type of transport. The step

of defining a new layer/network comprises the steps of: defining at least two Ethernet Access Points, namely Ethernet interfaces at the SDH/SONET network boundary where the Ethernet signals can access/leave the SDH/SONET network; defining a Link as a pair of Ethernet Access Points providing a point-to-point connection; for any pair of Ethernet Access Points, defining corresponding Circuits, namely all the possible routes connecting the pair of Access Points through the SDH/SONET network; and dividing each Circuit into Pipes, namely a sequence of smaller segments.

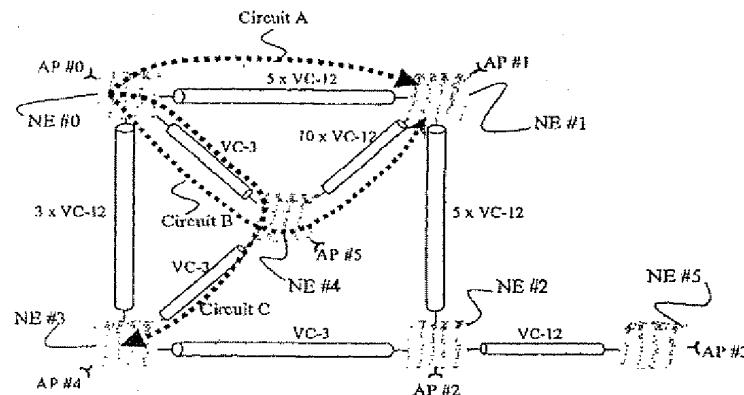
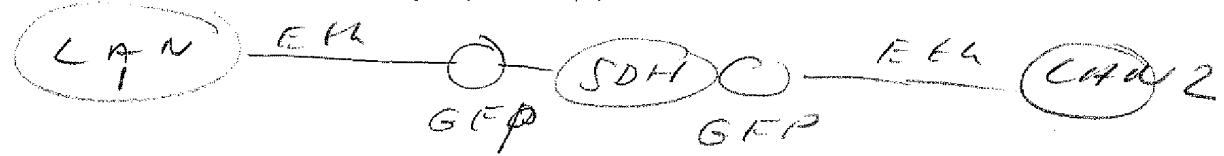


Fig. 1



Description

[0001] The present invention relates to the telecommunication field and in particular to a method and device for managing Ethernet frame traffic and transporting such a traffic over a transport SDH/SONET network.

[0002] As it is known, traffic generated by an Ethernet apparatus is characterized by discontinuities, namely there are periods with a more or less constant sending rate of Ethernet packets and periods during which a rather long time is provided between a received Ethernet frame and the next one. Such an unstable/inconstant traffic is generally termed "bursty". On the contrary, SDH or SONET traffic is characterized by a constant sending/receiving rate. In other words, any network element of a transport SDH/SONET network sends corresponding frames with a regular and constant rate. Furthermore, Ethernet frames do not have a fixed length/size but only a maximum size (1518 bytes).

[0003] It is easy to understand that these discrepancies result in a highly difficult interfacing of two technologies having different natures/characteristics.

[0004] An already available solution to the above problem allows the mapping of Ethernet frames into SDH/SONET Virtual Containers as a transparent tributary; all incoming bits are transported to the output interface with the related timing information (frequency for recovering the proper bit rate at the reception side). Within the SDH/SONET payload also the dead times between a received Ethernet frame and the following one are mapped.

[0005] Unfortunately, while such a solution could be considered easy to be implemented, the performances of this type of transport are the same as for the transport of a PDH tributary, namely rather low. This best prior solution is strictly tied to SDH architecture and it does not allow to provide new services or to achieve better performances with respect to the transport of a transparent PDH tributary. Furthermore, some bandwidth is wasted because it is used for transporting useless information.

[0006] In view of the above problems, the general object of the present invention is overcoming them in an efficient manner.

[0007] The main scope of the present invention is providing a method and device for an enhanced transport of Ethernet frame traffic over a transport SDH/SONET network.

[0008] The above and further objects of the present invention are obtained by a method and device according to claims 1 and 9, respectively. The present invention further comprises a network manager according to claim 8. Further advantageous features of the present invention are set forth in respective dependent claims. All the claims are intended as an integral part of the present description.

[0009] The basic idea of the proposed solution is to provide a complete new layer/network over the SDH/SONET network in order to manage the transport of Ethernet traffic over SDH/SONET network; this new layer/network uses the resources of SDH/SONET network in such a way as to optimize the provided services and the performances with reference to this specific type of transport.

[0010] The device according to the present invention is able to monitor continuously the Ethernet channel and to distinguish the nature of Carrier Event, so that the frames containing payload can be selected and the idle ones can be disregarded and not mapped in SDH virtual containers.

[0011] In order to increase the interfacing between Ethernet and SDH/SONET and better supporting all functionalities to be implemented, it has been decided to insert a further level of data encapsulation.

[0012] The present solution allows Ethernet clients to set-up their own Virtual Private Networks (based on point-to-point connections) by means of SDH/SONET network.

[0013] The present invention will become clear in view of the following detailed description, to be read having reference to the attached sheets of drawings, wherein:

- Fig. 1 shows the structure of a VPN and relating circuits;
- Fig. 2 is a schematic representation of encapsulation step of Ethernet frames into a SDH/SONET frame;
- Fig. 3 shows a basic scheme of insertion of GFP packets into C-4 containers; and
- Fig. 4 shows in better detail a selected Link with the two related Circuits.

[0014] The present invention is implemented by providing a complete new layer/network which is termed NETS (i.e. Network of Ethernet Transport over SDH/SONET). The NETS comprises basic elements to be defined herebelow.

[0015] The basic resources of NETS are the SDH/SONET Virtual Containers; the NETS uses these resources as basic pipelines to connect two Ethernet access points (point to point connection).

[0016] The NETS model provides different schemes of connection and management of these basic pipelines; by means of this new model it is possible to provide new services and to perform, in a better way, services already provided by SDH/SONET network.

[0017] The NETS model is based on five basic elements: Access Point, Link, Circuit, Pipe and Path.

[0018] An Access Point (AP) is an Ethernet Interface at the boundary of an SDH/SONET network; it is the point where the Ethernet traffic can access/leave the SDH/SONET network.

[0019] Fig. 1 depicts a simple example of network comprising six Network Elements (NE) with each network element having an Access Point; naturally, a Network Element can host more than one Access Point.

[0020] Only the Network Elements with capability to drop/insert Ethernet traffic are depicted; Network Elements that just manage SDH/SONET traffic are transparent and are not depicted.

[0021] A pair of Ethernet Access Points defines a point to point connection; this connection is named Link. For instance, with reference to Fig. 1, the pair AP #0 & AP #1 identifies a link; the couple AP #2 & AP #5 defines another link, and so on.

[0022] In case two access points are connected to a single link, such a link represents a Point-to-Point connection. Should an access point be connected to more than one link, such an access point is subject to a Point-to-multipoint connection but the various links are to be considered Point-to-Point connections. In case of multipoint connections, dispatching (sending different frames to several APs) and grooming (aggregations of frames coming from several APs) functionalities should be provided.

[0023] An SDH/SONET network could allow for the connection of two Access Points (i.e. to accomplish a Link) by means of different routes; every route is named Circuit. A Circuit is obtained by a Pipe concatenation and could be considered as a series connection of N Pipes.

[0024] The following Table 1 lists the possible routes for the Link identified by the Access Points AP #0 & AP #1 in Fig. 1.

Table 1

LINK	RELATED CIRCUITS
AP#0 - AP#1	NE #0 - NE #1
	NE #0 - NE#4 - NE #1
	NE #0 - NE #3 - NE #4 - NE #1
	NE #0 - NE #3 - NE #2 - NE #1
	NE #0 - NE #4 - NE #3 - NE#2 - NE #1

[0025] In principle, also route NE #0 - NE #4 - NE #3 - NE #0 - NE #1 is possible but it is not really significant because it is made up by a ring that leads to the starting point NE #0 and route #1.

[0026] Link AP #0 - AP #1 is accomplished by means of these five circuits; of course a subset of all possible Circuits could be selected to accomplish a Link.

[0027] In its turn, every Circuit/route that connects two Access Points can be divided into a sequence of smaller segments; every segment is named Pipe.

[0028] With reference to the previous list of Circuits, in Table 2 is the description of all the related Pipes.

Table 2

LINK	RELATED CIRCUITS	RELATED PIPES
AP #0 - AP #1	NE #0 - NE #1	NE #0 - NE #1
	NE #0 - NE#4 - NE #1	NE #0 - NE#4
	NE #0 - NE #3 - NE #4 - NE #1	NE#4 - NE #1
	NE #0 - NE #3 - NE #2 - NE #1	NE #0 - NE #3
	NE #0 - NE #4 - NE #3 - NE#2 - NE #1	NE #3 - NE #2

EP 1 339 198 A1

- [0029] As it is clear from the above Table 2, a Pipe can be shared among different Circuits.
 [0030] In its turn, every Pipe comprises one or more SDH/SONET Virtual Containers; this means that its capability is the sum of the capabilities of all the related Virtual Containers.
 [0031] Here below (see Table 3), with reference again to Figure 1, is a complete Link description with the Pipe composition.

Table 3

LINK	RELATED CIRCUITS	RELATED PIPES	PIPE COMPOSITION
AP #0 - AP #1	NE #0 - NE #1	NE #0 - NE #1	5 x VC - 12
	NE #0 - NE#4 - NE #1	NE #0 - NE#4	1 VC-3
		NE#4 - NE #1	10 x VC-12
	NE #0 - NE #3 - NE #4 - NE #1	NE #0 - NE #3	3 x VC-12
		NE #3 - NE #4	1 VC-3
		NE #4 - NE #1	10 x VC-12
	NE #0 - NE #3 - NE #2 - NE #1	NE #0 - NE #3	3xVC-12
		NE #3 - NE #2	1 VC-3
		NE #2 - NE #1	5 x VC-12
	NE #0 - NE #4 - NE #3 - NE#2 - NE #1	NE #0 - NE #4	1 VC-3
		NE #4 - NE #3	1 VC-3
		NE #3 - NE#2	1 Vc-3
		NE#2 - NE #1	5 x VC-12

- [0032] The basic pipeline is the Virtual Container that connects two Network Elements; it is named Path.
 [0033] In order to improve the interfacing between Ethernet and SDH/SONET, and for better supporting all the functionalities that have been selected, it has been decided to insert a further level of data encapsulation. The protocol used for this intermediate level is GFP (Generic Frame Procedure). This technique has been selected because GFP packet has been properly defined for allowing the mapping of generic data in transport frames having variable-length payload and based on octet alignment, such as SDH/SONET and OTN (Optical Transport Network). Fig. 3 shows a basic scheme of data encapsulation according to the present invention.
 [0034] From the figure one can infer that the first encapsulation stage has a 1:1 rate (an Ethernet message is inserted in a GFP message). As far as the mapping of GFP packets into SDH/SONET Virtual Containers is concerned, such a mapping is accomplished by a method similar to the one used for mapping ATM frames. In order to facilitate the comprehension of such a procedure, it is possible to consider the GFP packets as a byte stream to be inserted into Virtual Containers. The number of GFP packets that could be inserted into a VC is variable and depends on two fundamental factors: the VC capacity and the GFP dimension which in turn depends on the Ethernet frame to be transported.
 [0035] According to these two factors, it is possible that a single GFP packet is mapped into two or more Virtual containers or that a plurality of GFP packets is inserted into a single VC.
 [0036] In order to recover the GFP frames, it is necessary to perform an alignment to PLI field, which in turn indicates the length of payload data field. The PLI (PDU Length Indication) field will be disclosed below. Some examples of GFP-to-VC-x mapping are also reported below.
 [0037] In PLI field there are provided two octets. They are the binary number representing the number of octets contained into the payload area of the packet itself.
 [0038] The monitor information for managing the network according to the present invention are exchanged through the use of proper Scout Message GFP packets. They are control packets that are always sent before sending single packets containing Ethernet frames (In principle, a proper scout message alone is sent in case there are no messages to be transported).
 [0039] Four different scout messages are provided according to the present invention. In principle, each different scout message covers a portion of the information that is requested by the network that is realized in accordance with the present invention. Conveniently, each Scout Message comprises a SMT (Scout Message Type) field for clearly indicating the type of scout message (for instance: STM for Path Status Message = 00; STM for Complete Status

Message = 01; STM for Complete Status and Ethernet Message Info = 02 and STM for Complete Status and Delay Message = 03).

[0040] The first and simpler Scout Message that is provided is termed *Path Status Message*. It is able to report only the information relating to the path operation condition.

5 [0041] The second Scout Message, termed *Complete Status Message*, contains all the information required for calculating an estimation of delays of data propagation through the virtual private network and the information relating to the operation status of Links, Circuits and Paths.

10 [0042] The third Scout Message, termed *Complete Status and Ethernet Message Info*, is put before the GFP message containing the Ethernet frame. The object of such a packet is to transport all the information relating to the operation status, those information required for calculating the transit time of packets and those information relating to the Ethernet message that is transported.

15 [0043] The fourth Scout Message, termed *Complete Status and Delay Message*, is very similar to the third one but the difference is in that it is sent without a GFP packet containing an Ethernet frame. The object of such a Scout Message is to provide useful indications about transit times of Circuits of a Link even if there are no Ethernet frames to be sent. Just for this reason, all the fields containing information relating to the data message are not significant; only the fields transporting the Path status, the Circuit status, Link status and their delays are considered as valid.

20 [0044] There now follow some examples of mapping Ethernet frames into SDH structures. The very same concepts are equally applicable to SONET structures.

[0045] In order to better understand the mapping procedure according to the present invention, the following basic notions and new concepts are set forth.

Capacity C-4: 2340 bytes

Capacity C-12: 136 bytes

Capacity C-3: 756 bytes

25 Max size of Ethernet frames: 1518 bytes + 8 bytes (preamble + SDF)

Size of additional GFP fields for an Ethernet frame: 12 bytes

Size of Scout Message *Complete Status and Ethernet Message Info*: 20 bytes

Size of Scout Message *Complete Status Message*: 18 bytes

30 [0046] First Example: mapping a 1526-byte Ethernet frame, together with an accompanying Scout Message, into i) a VC-4, ii) a VC-3 or iii) VC-12. The accompanying Scout Message is *Complete Status and Ethernet Message Info*. Within the GFP packet encapsulating the Ethernet Frame, the Preamble and SFD (Start Frame Delimiter) fields of the Ethernet frame should not be included; thus, in the Payload Data field of GFP, 1518 bytes are inserted. As the Scout Message size is 20 bytes, 1550 bytes are inserted into the SDH frame payload.

35 i) Mapping into a VC-4: a C-4 container has 2340 bytes and thus it is possible to insert only the Ethernet message together with the respective Scout Message ($2340 \text{ bytes} / 1550 \text{ bytes} = 1,509$). In case two Ethernet frames should be transported, it is clear that the second Ethernet frame and its Scout Message can be only partially inserted in the first container while the remaining bytes should be mapped into the second C-4 (see Fig. 3).
 40 ii) Mapping into a C-3: a C-3 container has 756 bytes and thus three C-3 containers should be used for inserting an Ethernet frame. Two containers will be completely filled while the third one will be only partially used.
 iii) Mapping into a C-12: a C-12 has 136 bytes and thus twelve C-12 containers should be used for inserting an Ethernet frame. Eleven containers will be completely filled while the last one will be only partially used.

45 [0047] Second Example: mapping a 18-bytes *Complete Status Message* into iv) C-4, v) C-3 and vi) C-12.

iv) Mapping into a VC-4: a C-4 container has 2340 bytes and thus it is possible to insert 130 *Complete Status Messages*.

v) Mapping into a C-3: a C-3 container has 756 bytes and thus it is possible to insert 42 *Complete Status Messages*.
 vi) Mapping into a C-12: a C-12 has 136 bytes and thus seven *Complete Status Messages* could be inserted, with a further *Complete Status Message* being partially inserted.

[0048] Naturally, the above calculations are purely theoretic, as the size of an Ethernet frame does not always correspond to the admissible maximum size.

55 [0049] A point-to-point connection of two Access Points is accomplished by the following steps:

- a) Defining a Link;
- b) Selecting one or more Circuit(s);

- c) Defining the size of the related Pipes; and
d) Selecting the related Paths.

As now the point-to-point connection is fully defined, the simple transport of an Ethernet frame is performed by the following steps:

- e) Receiving the Ethernet frame at the Ethernet Interface of one Access Point;
f) Routing the Ethernet frame to at least one link;
g) Selecting one of the available Circuits;
h) Selecting one available Path of the (first) Pipe of the selected Circuit;
i) Encapsulating the Ethernet frame into the selected available Path (i.e. Virtual Container);
j) Transporting the Ethernet frame up to the next Network Element (namely, up to the end of the Pipe); and
k) Extracting the Ethernet frame from the Virtual Container.

[0050] In case the selected Circuit comprises further Pipes, steps h) to k) are repeated in every intermediate Network Element until the last Network Element is reached. Furthermore

l) Inserting the extracted Ethernet frame in a re-ordered queue. In fact, due to the possible skew among different Circuits or among the different Paths of a Pipe, the order of the messages received at the destination Network Element could be different from the order of the messages at the starting Access Point, so a re-ordering action is required. Should just one Circuit and simple Paths be used, step k) is not required.

m) Finally, providing the Ethernet frame to the Ethernet interface of the destination AP.

[0051] The above is just an example of simple transport for showing the versatility of the additional network/layer according to the present invention. Herebelow is a detailed description of an Ethernet hitless protection performed through the present invention. In principle, SDH/SONET networks already provide different types of protection (for instance SNCP or MS-SPRING) that can be applied to Ethernet frames as Ethernet frames are encapsulated into SDH/SONET Virtual Containers. The advantage of the Ethernet hitless protection mechanism according to the present invention is that it is performed at the lower possible level, it could be easily hardware implemented and provides hitless performances.

[0052] With reference to Fig. 1, let consider the point-to-point connection (Link AP #0 - AP #1) which is identified by the pair of Access Points AP #0 and AP #1. Different routes made up by SDH/SONET Virtual Containers can connect the two Access Points AP #0 and AP #1; for instance, two of them could be Circuit A and Circuit B. Circuit A is the direct route comprising five VC-12; Circuit B is a route comprising a sequence of one VC-3 and ten VC-12 with an intermediate node (NE #4). In principle, several other routes can connect the two Access Points AP #0 and AP #1 but for the aim of this example and for clarity reasons just two of them will suffice.

[0053] Fig. 5 highlights the selected Link with the two related Circuits.

[0054] The basic idea of the proposed solution comprises the following steps:

- Every time an Ethernet frame is received at AP #0 of NE #0 it is transmitted along both Circuit A and Circuit B. Clearly, the transmission along two different routes results in a protection service.
- As a consequence of the previous step, NE #1 receives the same Ethernet frame twice; as a rule, it will accept the frame received from the faster Circuit and discharges the second one. Let consider Circuit B is faster than Circuit A; Ethernet frames received from Circuit B are selected while the frames received from Circuit A are discharged.
- In case of failure of Circuit B, NE #1 just receives frames from Circuit A and of course accepts them; the protection is accomplished by changing the selection of the Circuit from B to A.

[0055] The protection is hitless because the selection is frame-based and the sequence of Ethernet frames is maintained.

[0056] Here is a more detailed description of the proposed Ethernet hitless protection solution, having further reference to Figs. 1 and 4.

[0057] Every time an Ethernet frame is received at AP #0 of NE #0 it is stored in a queue of incoming messages.

[0058] Each frame is labeled in order to recover the exact frame sequence at the ending point. For instance, a sequence of labeled received frames could be FR₁, FR₂, FR₃, ... FR_n.

[0059] The received frame is transmitted by two separate transmitters, TXA and TXB, to Circuit A and Circuit B, respectively. Different types of Virtual Containers (VC-12 for Circuit A and VC-3/VC-12 for Circuit B) perform the transport of a frame along different routes (direct for Circuit A, with an intermediate node for Circuit B). At the receiving node two different receivers are provided (RXA, RXB).

[0060] At intermediate node (NE #4) of Circuit B, the Ethernet frames are stored in an intermediate node queue. In

principle, the presence of an intermediate node in a Circuit results in a delay in the frame transmission. In the present case, Circuit B is faster than Circuit A because the capability of the two pipes (one VC-3 and ten VC-12) of Circuit B is higher notwithstanding the presence of the intermediate node (NE #4).

[0061] This means that a frame (for instance FR₁) from Circuit B is received at NE #1 before the same frame is received from Circuit A; NE #1 just selects the first received one and stores it in a queue of outgoing Ethernet frames that are transmitted at AP #1.

[0062] Until both the Circuits are active, only the frames received from the faster one are selected and stored in the queue of outgoing frames; of course the frame sequence is maintained (possibly it is recovered through the label).

[0063] If a failure occurs on Circuit B, after the transmission of frame FR₃, NE #1 receives frame FR₄ only once (only from Circuit A).

[0064] NE #1 selects frame FR₄ from Circuit A because it is the first received one and stores it in the queue; the same for the following frames until Circuit B is restored.

[0065] Frame FR₄ and the following ones are not lost and the frame sequence is maintained; this means that a hitless protection was successfully performed.

[0066] When Circuit B is restored, NE #1 changes again its selection into a failure-free hitless mode.

[0067] It is important to note that the Circuit selection is frame-by-frame based and it is not related neither to the last selection nor to the reception status of the previous or the following frames. In other words, when a frame labeled as FR_{n+1} is received, the Circuit selection is performed independently of the last selection and also independently whether frame FR_n or FR_{n+2} has already been received or not.

[0068] Anyway, the reception of frame FR_{n+1} before frame FR_n could occur when the skew between the two Circuits is equal to or longer than the time required for a frame transport or the faster Circuit is restored after a failure.

[0069] It could happen that the transport of frame FR_n is performed by Circuit A only because of the failure of Circuit B but the transport of frame FR_{n+1} is performed by both Circuits because Circuit B has been restored. Due to the skew between the two Circuits, NE #1 receives frame FR_{n+1} from Circuit B before frame FR_n from Circuit A. Frame FR_{n+1} is stored in the queue and the related selection is Circuit B; also frame FR_n is stored in the queue but the related selection is Circuit A. Of course, frame FR_{n+1} cannot be provided to AP #1 until the reception of frame FR_n.

[0070] The independent frame-by-frame selection of the Circuit is important not only when a Circuit failure occurs/disappears but also when both the Circuits are active.

[0071] With reference again to Fig. 1, let consider a further Link, the Link AP #0 - AP #4; one Circuit of this Link (Circuit C) could be made up by a VC-3 between NE #0 and NE #4 and a VC-3 between NE #4 and NE #3. It is realized that the first VC-3 is the same VC-3 used for Circuit B of Link AP #0 - AP #1, i.e. it is a shared resource.

[0072] The above results in a transport delay of a frame along Circuit B that also depends on the traffic of the other Circuit (Circuit C) and that can dynamically change.

[0073] In case Circuits A, B and C are active and the transport delay along one or both Circuits dynamically changes for the shared resources, the fastest Circuit can dynamically move from Circuit A to Circuit B and vice versa. Thus, the Circuit selection can change also without the occurrence of any failures.

[0074] The present invention can be implemented both in hardware and software. Advantageously it is hardware implemented through a SDH/SONET network comprising network elements (for instance ADMs and Cross-Connects) and fiber connections. In particular, the new layer according to the invention is provided by a network manager managing the physical network at an high level. Furthermore, within the network elements (or at least a part thereof) at least an additional board is provided. Each additional board comprises at least one Ethernet interface, namely an Access Point. Generally, a number of Access Points are provided in each additional board.

[0075] According to a preferred embodiment of the present invention, each additional board comprises FPGA means (two FPGAs, namely two Field Programmable Gate Arrays), memory means and integrated Circuit means (two ASICs).

The network manager provides some information to the additional board (particularly to the FPGAs) comprising which AP should be used, the bit rate of the Ethernet flow (10 or 100 Mb/s) and the SDH/SONET resources to be used for transporting the Ethernet signal. Furthermore, the FPGAs perform several additional tasks such as filling/emptying the virtual containers.

[0076] The memory means comprise several memories, namely a data memory, an external memory with routing information, a link memory for storing information about each link and a circuit memory for containing the circuit queues and tables.

[0077] A further advantage that is provided by the present invention is that each GFP packet for Ethernet frames comprises a Core Header Error Check field containing a CRC error correction code for protecting the integrity of GFP packet core header. The CRC error correction code according to the present invention is able to correct a single error and to detect any possible further errors. Thus, the advantage is that only the errored Ethernet frames could be discarded when a SDH/SONET frame is received, the error-free frames could be advantageously kept. This is in contrast with the error correction code mechanisms that are provided for correcting errors in the whole SDH/SONET frame.

[0078] A still further advantage of the present invention is that it could be applied to any network topology, namely

linear, meshed, ring, tree...

Claims

5 1. A method for handling Ethernet frame signals in a SDH/SONET network, the SDH/SONET network comprising network elements or nodes and fiber connections connecting the network elements, the method being characterized by the step of defining a new layer/network (NETS) over the SDH/SONET network in order to manage the Ethernet signals over the SDH/SONET network, the new layer/network (NETS) using the resources of SDH/SONET network in such a way as to optimize the provided services and the performances with reference to this specific type of transport, wherein the step of defining a new layer/network (NETS) comprises the steps of:

10 defining at least two Ethernet Access Points (AP), namely Ethernet interfaces at the SDH/SONET network boundary where the Ethernet signals can access/leave the SDH/SONET network;
 15 defining a Link as a pair of Ethernet Access Points providing a point-to-point connection;
 for any pair of Ethernet Access Points, defining corresponding Circuits, namely all the possible routes connecting the pair of Access Points through the SDH/SONET network; and
 dividing each Circuit into Pipes, namely a sequence of smaller segments.

20 2. Method according to claim 1, characterized in that it further comprises the steps of:

25 e) receiving the Ethernet frame signals at the Ethernet interface of one Access Point (AP);
 f) routing the Ethernet frame signals to at least one link;
 g) selecting one of the available Circuits;
 h) selecting one available Path of the first Pipe of the selected Circuit;
 i) encapsulating the Ethernet frame signals into at least one Virtual Container of the selected available Path;
 j) transporting the Ethernet frame signals to the next Network Element namely, up to the end of the Pipe; and
 k) at the Ethernet interface of the destination Access Point, extracting the Ethernet frame signals from the at least one Virtual Container.

30 3. Method according to claim 2, characterized by comprising the additional step of inserting the extracted Ethernet frame in a re-ordered queue.

35 4. Method according to claim 2, characterized in that, in case the selected Circuit comprises further Pipes, steps h) to i) are repeated in any intermediate Network Element until the last Network Element is reached.

40 5. Method according to claim 2, characterized in that step i) comprises the step of mapping the Ethernet frame signals into GFP frames and mapping the GFP frames into Virtual Containers.

45 6. Method according to claim 2, characterized in that step i) comprises the step of providing a scout message containing monitor information for managing the new layer/network (NETS).

50 7. Method according to claim 2, characterized in that step j) comprises the step of transporting the same Ethernet frame signals through at least two different Circuits and step k) comprises the step of performing a hitless frame-based switch in order to select frames from the faster Circuit or from the non failed Circuit.

55 8. Network manager for managing a SDH/SONET network and handling Ethernet frame signals therein, the SDH/SONET network comprising network elements or nodes and fiber connections connecting the network elements, the manager being characterized by providing a new layer/network (NETS) over the SDH/SONET network in order to manage the Ethernet signals over the SDH/SONET network, the new layer/network (NETS) using the resources of SDH network in such a way as to optimize the provided services and the performances with reference to this specific type of transport, wherein the step of providing a new layer/network (NETS) comprises:

defining at least two Ethernet Access Points (AP), namely Ethernet interfaces at the SDH/SONET network boundary where the Ethernet signals can access/leave the SDH/SONET network;
 defining a Link as a pair of Ethernet Access Points providing a point-to-point connection;
 for any pair of Ethernet Access Points, defining corresponding Circuits, namely all the possible routes connecting the pair of Access Points through the SDH/SONET network; and

dividing each Circuit into Pipes, namely a sequence of smaller segments.

9. Device for handling Ethernet frame signals in a SDH/SONET network, the SDH/SONET network comprising network elements or nodes and fiber connections connecting the network elements, **characterised in that** it comprises:

5 at least one access point (AP) interface for receiving Ethernet frame signals;
means for routing the received Ethernet frame signals to at least one link, wherein a link is a pair of Ethernet
Access Points providing a point-to-point connection in the network;
10 means for selecting an available Circuit, wherein a Circuit is one of the possible routes connecting the pair of
Access Points through the network; and
means for encapsulating the Ethernet frame signals into at least one Virtual Container of the selected available
Path.

- 15 10. Device according to claim 9, **characterised in that** it further comprises means for extracting the Ethernet frame
signals, transported through the network, from the at least one Virtual Container.

11. Device according to claim 9, **characterised in that** it further comprises memory means for storing the Ethernet
frame signals in signal queues.

- 20 12. Device according to any claims 9-11, **characterized in that** it comprises means for handling information received
by a network manager, such information comprising indications of access points to be used, bit rate of incoming
Ethernet frames and resources to be associated to a certain access point.

25

30

35

40

45

50

55

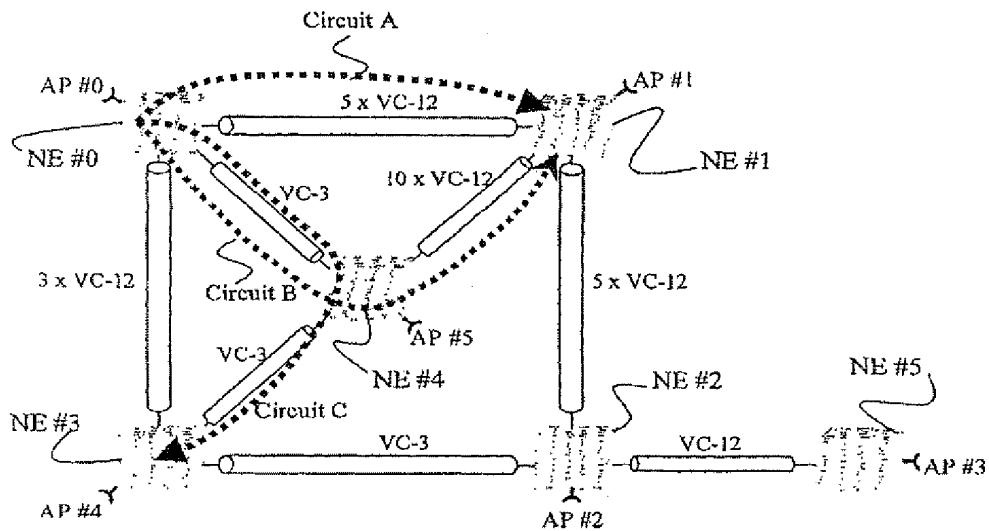


Fig. 1

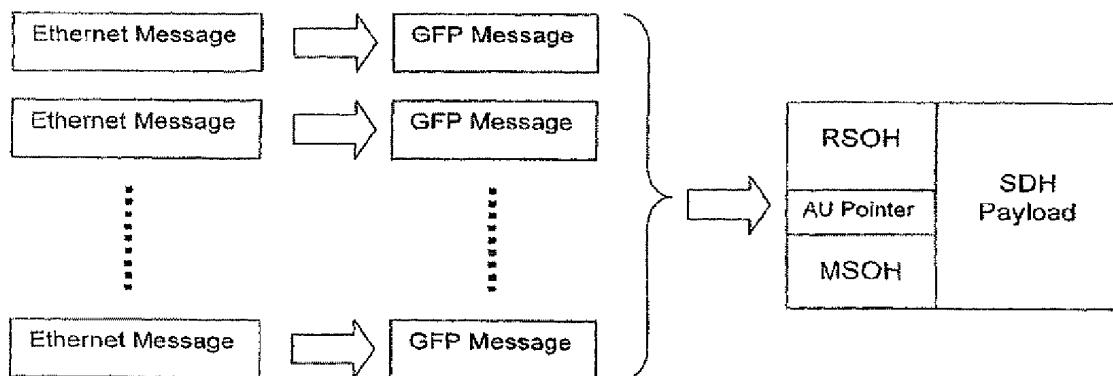


Fig. 2

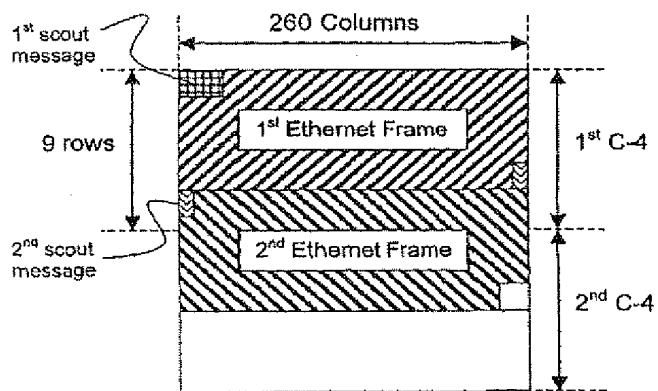


Fig. 3

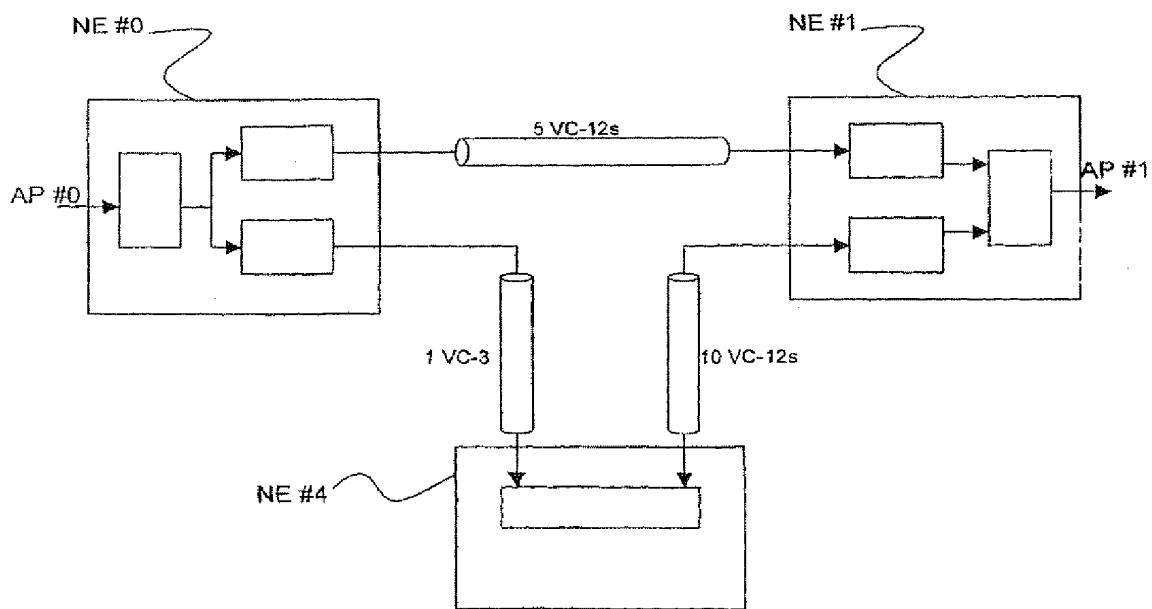


Fig. 4

European Patent
Office

EUROPEAN SEARCH REPORT

Application Number

EP 02 29 0445

DOCUMENTS CONSIDERED TO BE RELEVANT			CLASSIFICATION OF THE APPLICATION (H.O.L.)
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	
Y	WO 01 15363 A (MARCONI COMM INC) 1 March 2001 (2001-03-01) * abstract * * page 2, line 24 - line 29 * * page 4, line 1 - line 20 * * page 6, line 1 - line 10 * * page 8, line 5 - page 9, line 12 * * page 9, line 13 - line 28 * * figures 1,2,6 *	1-12	H04L29/06 H04L12/46 H04J3/16
Y	US 2001/043603 A1 (YU SHAOHUA) 22 November 2001 (2001-11-22) * abstract * * page 2, paragraph 19 - paragraph 24 * * page 6, paragraph 87 - paragraph 89 * * page 8, paragraphs 143,153 * * page 9, paragraph 183 * * figures 1,2,15 *	1-12	
A	US 6 122 281 A (AUGUSTA STEPHEN D ET AL) 19 September 2000 (2000-09-19) * figure 9 *	2	TECHNICAL FIELDS SEARCHED (H.O.L.)
A	EP 0 982 900 A (NORTEL NETWORKS CORP) 1 March 2000 (2000-03-01) * abstract * * page 6, line 16 - line 23 *	2	
A	EP 0 924 901 A (NORTHERN TELECOM LTD) 23 June 1999 (1999-06-23) * abstract * * column 8, line 54 - line 55 *	3	
		-/-	
The present search report has been drawn up for all claims			
Place of search	Date of completion of the search	Examiner	
THE HAGUE	3 July 2002	Bernardini, A	
CATEGORY OF CITED DOCUMENTS			
X : particularly relevant; if taken alone	T : theory or principle underlying the invention		
Y : particularly relevant if combined with another document of the same category	E : earlier patent document, but published on, or after the filing date		
A : technological background	D : document cited in the application		
O : non-written disclosure	L : document cited for other reasons		
P : intermediate document	Q : member of the same patent family, corresponding document		



European Patent
Office

EUROPEAN SEARCH REPORT

Application Number
EP 02 29 0445

DOCUMENTS CONSIDERED TO BE RELEVANT			
Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int.Cl.)
A	<p>ARMSTRONG T ET AL: "GFP FOR ETHERNET" CONTRIBUTION TO T1 STANDARDS PROJECT, XX, XX, 10 July 2000 (2000-07-10), pages 1-6, XP000949000 * abstract * * page 2 - page 4 *</p> <p>-----</p> <p>A "SERIES G: TRANSMISSION SYSTEMS AND MEDIA, DIGITAL SYSTEMS AND NETWORKS TYPES AND CHARACTERISTICS OF SDH NETWORK PROTECTION ARCHITECTURES" ITU-T RECOMMENDATION G.841, XX, XX, October 1998 (1998-10), page COMPLETE XP000955617 * paragraph '08.1! *</p> <p>-----</p>	5 7	
			TECHNICAL FIELDS SEARCHED (Int.Cl.)
<p>The present search report has been drawn up for all claims</p>			
Place of search	Date of completion of the search	Examiner	
THE HAGUE	3 July 2002	Bernardini, A	
CATEGORY OF CITED DOCUMENTS		T : theory or principle underlying the invention E : earlier patent document, but published on, or after the filing date D : document cited in the application L : document cited for other reasons S : member of the same patent family, corresponding document	
X : particularly relevant if taken alone Y : particularly relevant if combined with another document of the same category A : technological background O : non-written disclosure P : Intermediate document			

**ANNEX TO THE EUROPEAN SEARCH REPORT
ON EUROPEAN PATENT APPLICATION NO.**

EP 02 29 0445

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report. The members are as contained in the European Patent Office EDP file on. The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

03-07-2002

Patent document cited in search report		Publication date		Patent family member(s)	Publication date
WO 0115363	A	01-03-2001	AU EP WO	6616300 A 1206858 A1 0115363 A1	19-03-2001 22-05-2002 01-03-2001
US 2001043603	A1	22-11-2001	CN AU WO CN	1250294 A 6146200 A 0108356 A1 1316146 T	12-04-2000 13-02-2001 01-02-2001 03-10-2001
US 6122281	A	19-09-2000	AU AU CA DE EP WO	707746 B2 3732897 A 2260054 A1 69711053 D1 0916213 A1 9804072 A1	15-07-1999 10-02-1998 29-01-1998 18-04-2002 19-05-1999 29-01-1998
EP 0982900	A	01-03-2000	EP	0982900 A2	01-03-2000
EP 0924901	A	23-06-1999	US EP	6222848 B1 0924901 A2	24-04-2001 23-06-1999